

To Share or Not to Share: a case for MPI in Shared-Memory

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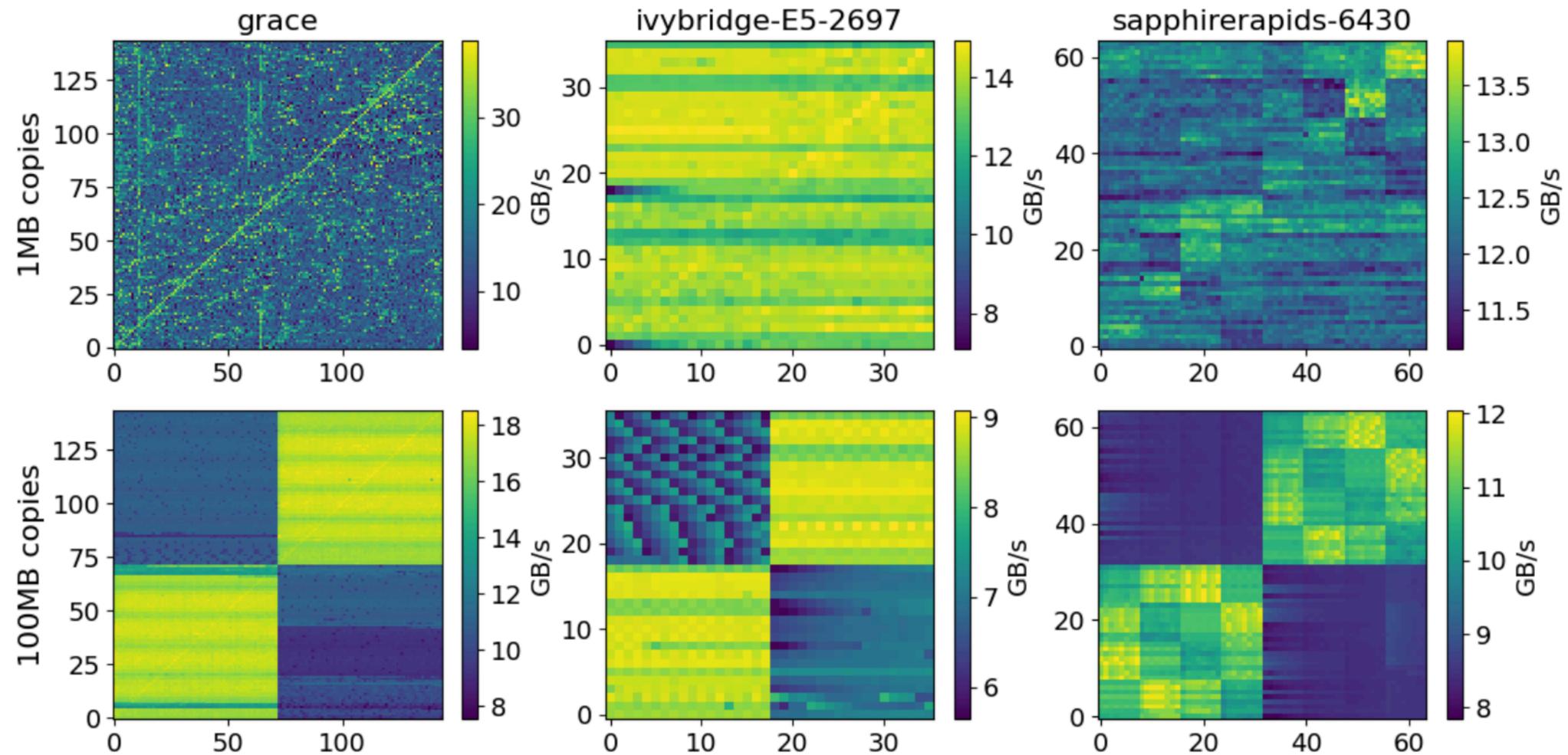
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91680 Bruyères-le-Châtel, France

Discussing how Shared-Mem can Benefit to MPI

- The Obvious Starting Point
- MPI and Compute Topology
- Horizontalization of Compute
- Current Support of MPI in SHMEM
- Sessions at the Rescue

The Obvious Starting Point

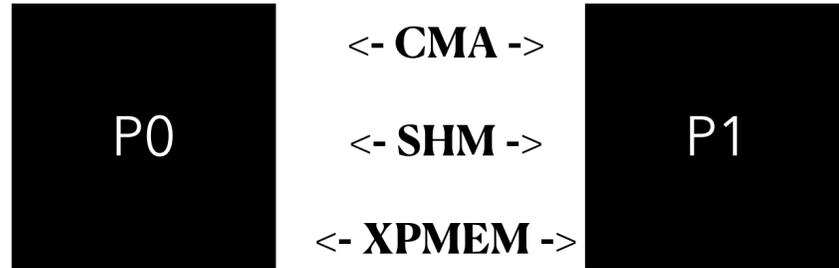
Nodes Have a Topology



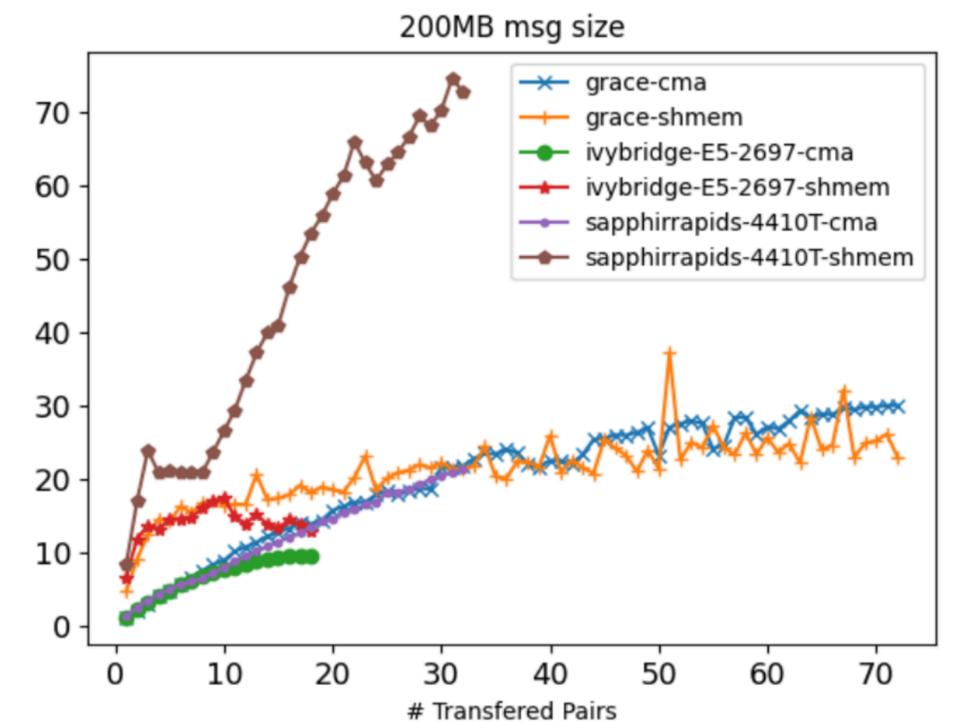
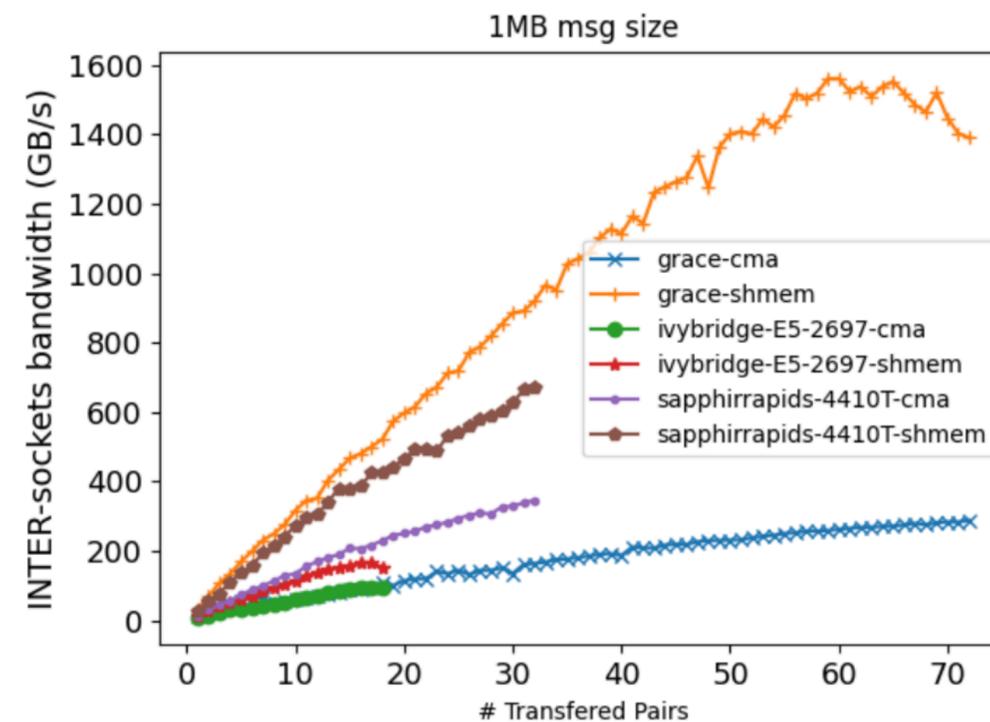
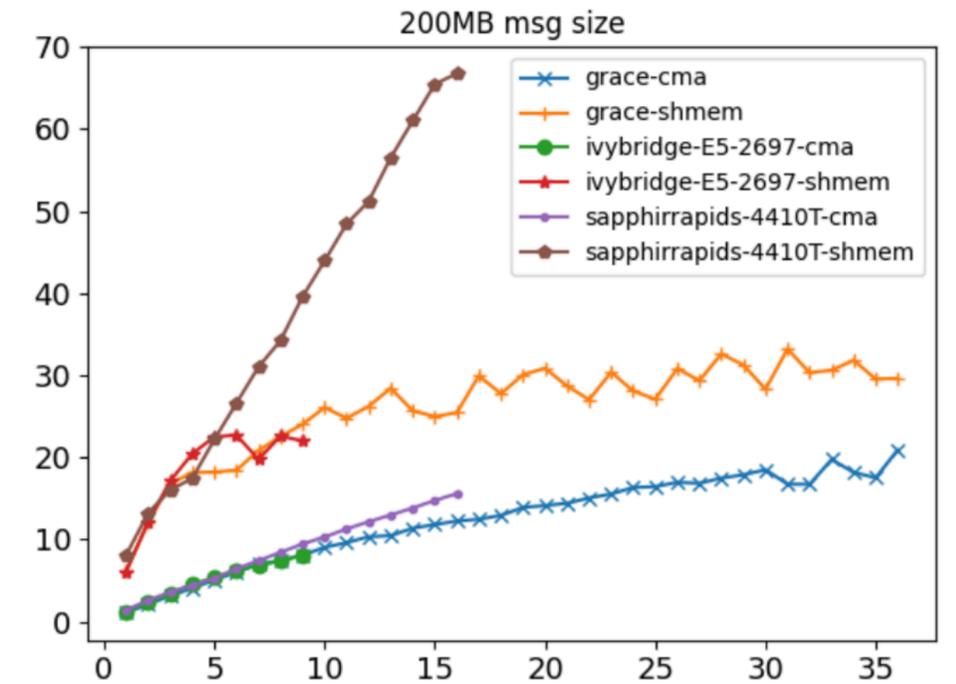
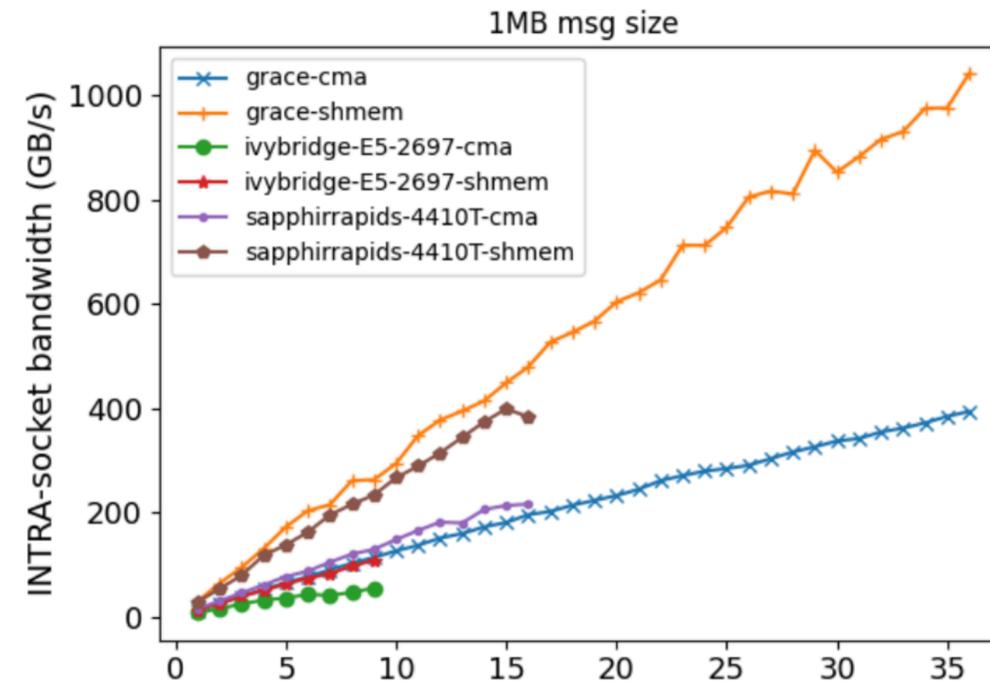
Core to core bandwidth

Performance in Shared-Memory

v- SHMEM -v



- Shared-memory transfers give the best performance compared to CMA when doing trivial benchmarks
- Note XPMEM was not deployed in test environment



MPI and Compute Topology

Seeing MPI as the Embedding Model



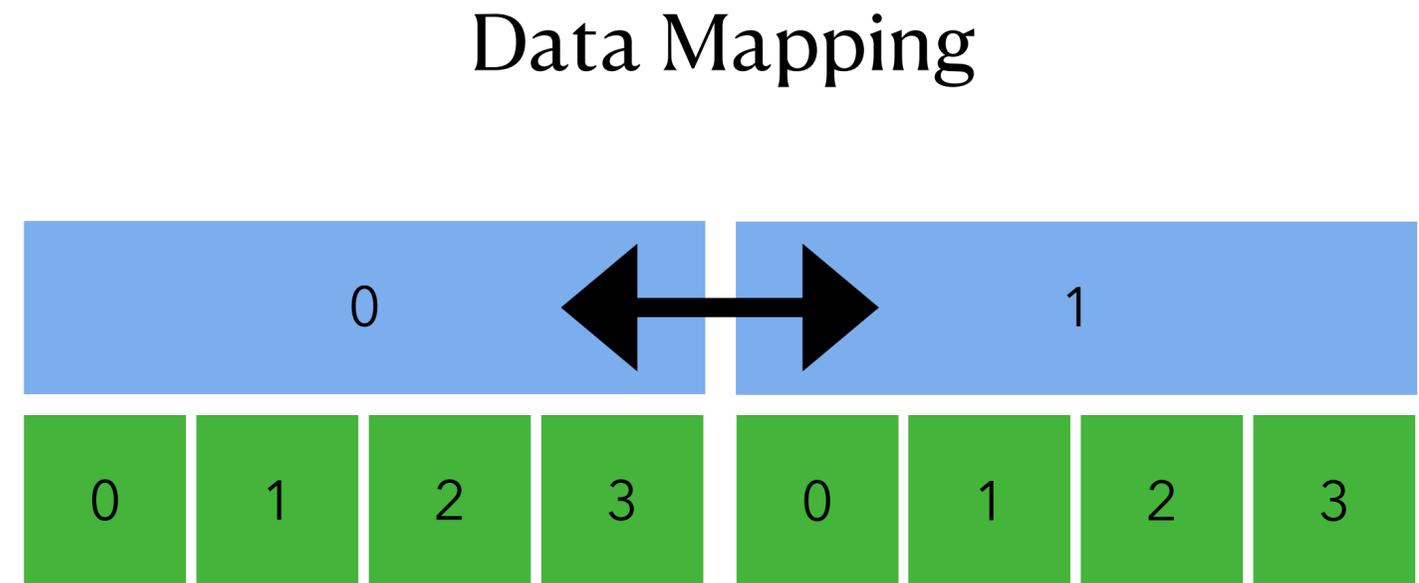
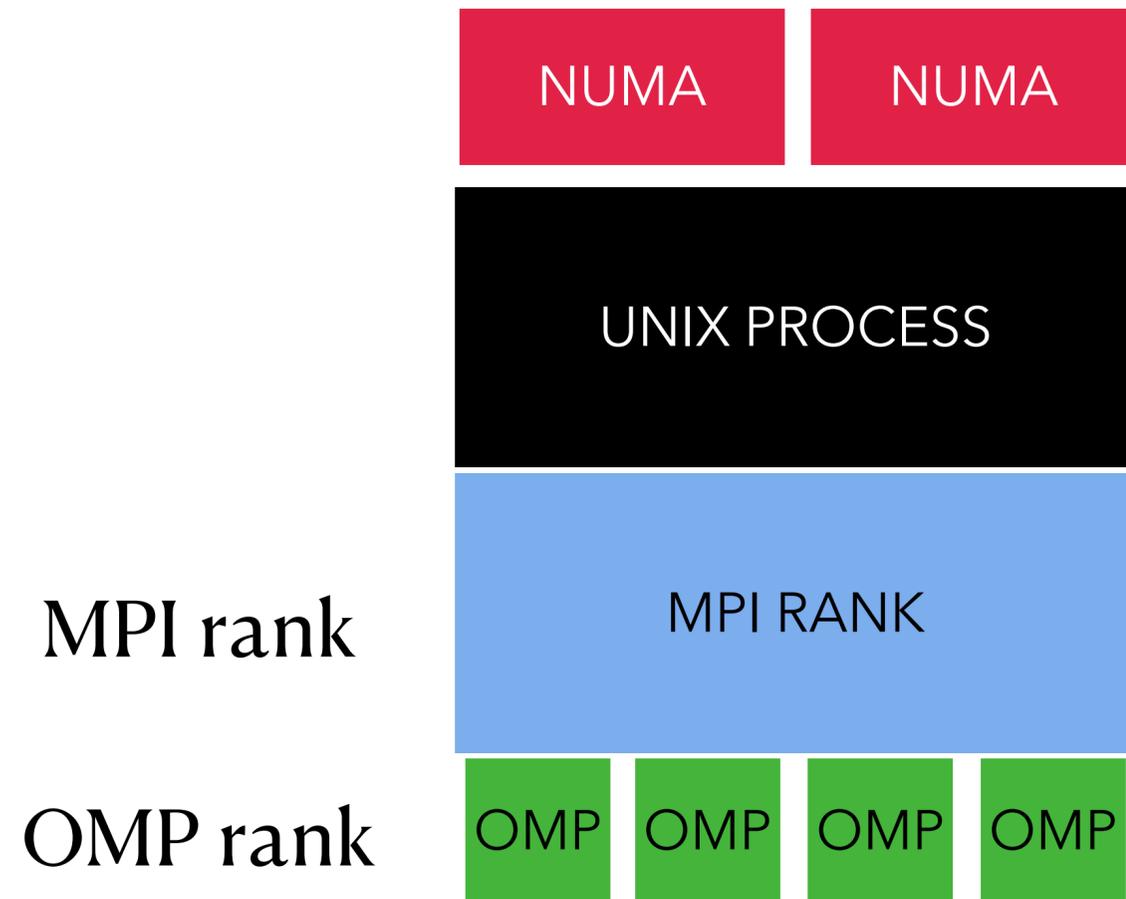
MPI is in charge of **Mapping** and **Addressing** the distributed ranks.

Seeing MPI as the Embedding Model



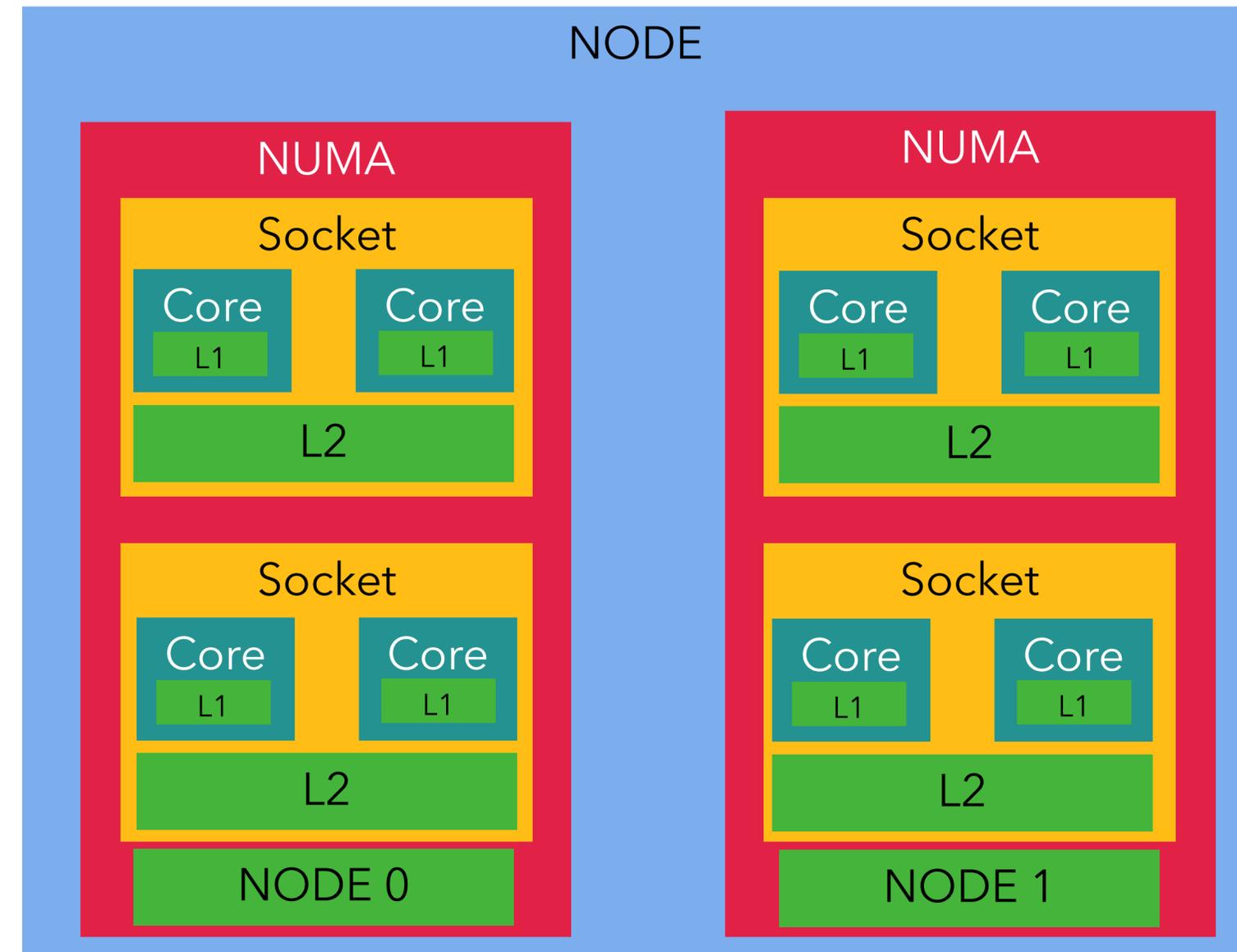
OMP Pinning is generally inherited from MPI

Seeing MPI as the Embedding Model



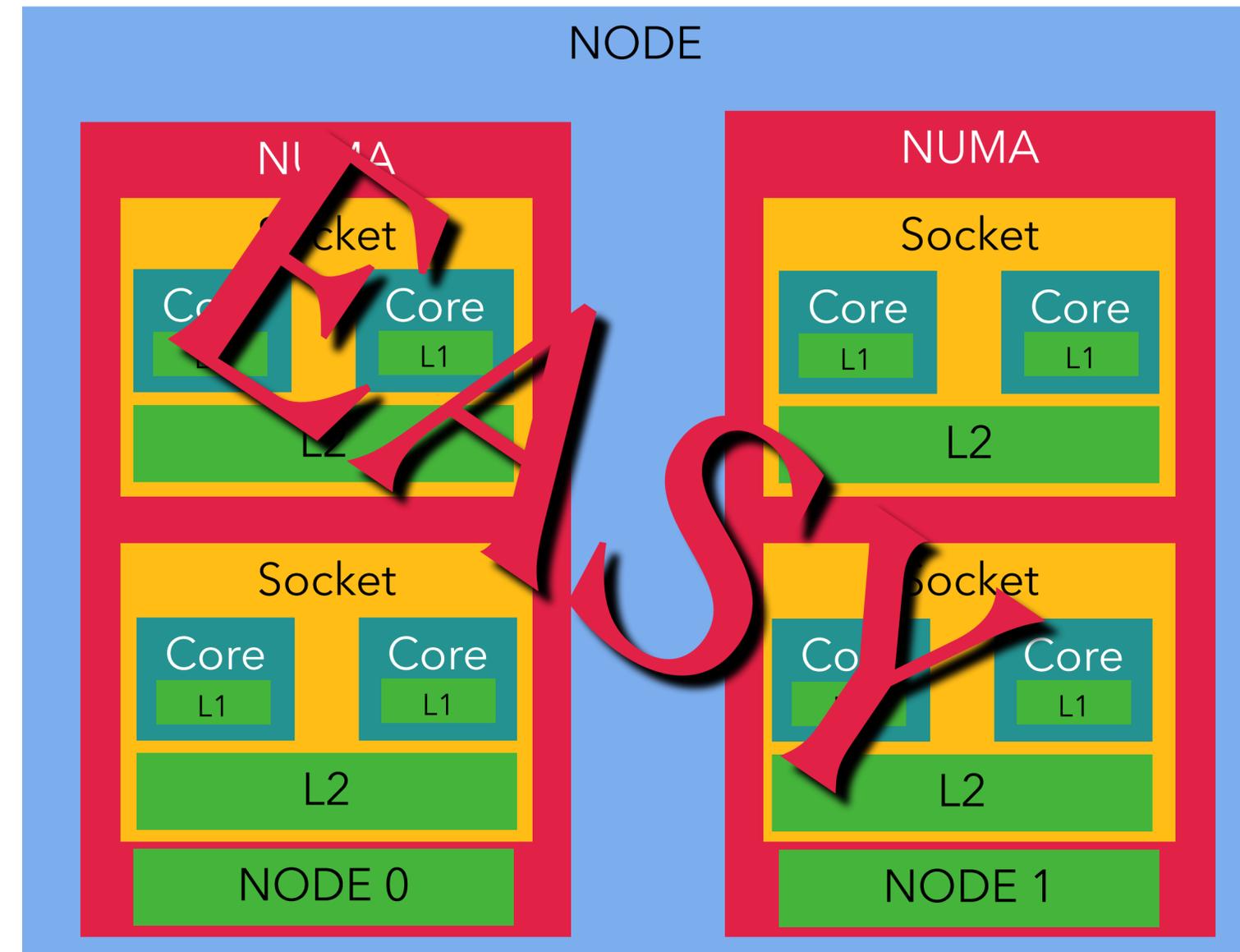
Parallel Programs' Spatialization

- With differentiated hardware the program does not only live in **time** but also in **space**
- Nesting programming models creates a hierarchy for locality with different indexing (MPI rank, OMP rank, Cuda indexes)
- Data must follow these indexes as they are used to define computational split
- Locality is increasingly a key performance aspect of current systems



Parallel Programs' Spatialization

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Parallel Programs' Spatialization



Nvidia GH100

8 Compute Graphic Cluster (GPC); 72 Texture Processing Clusters (TPC); 144 Streaming Multiprocessor (SM); 18 432 Cuda Cores

Parallel Programs' Spatialization

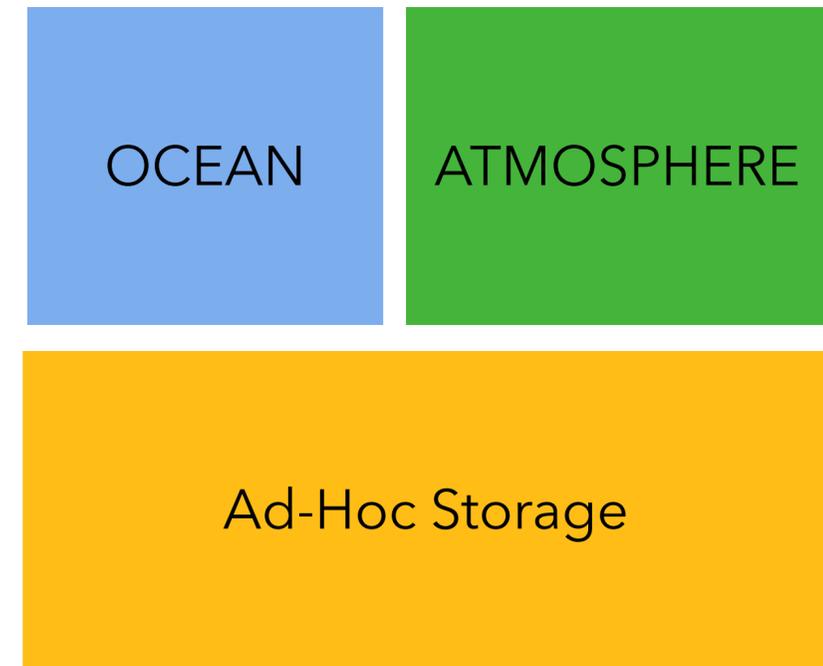


Current architectures also features a regular CPU over coherent memory

Horizontalization of Compute

On Software Differentiation

- Maintaining programs in an SPMD model is increasingly difficult:
 - Hardware is not symmetrical
 - Leads to nested states and programming models
- There is a need for **software differentiation**:
 - Collaborating programs require communication
 - Horizontalization of compute and not only data (including programs)

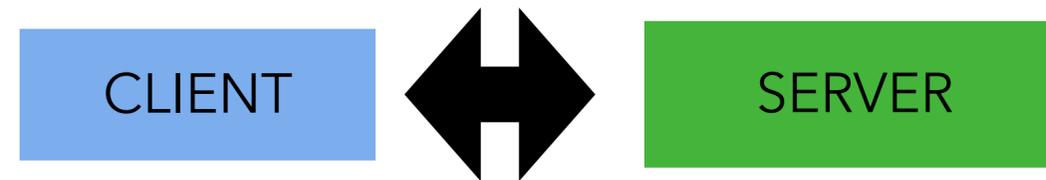


On Software Differentiation

MPI is not bulk synchronous, but ...



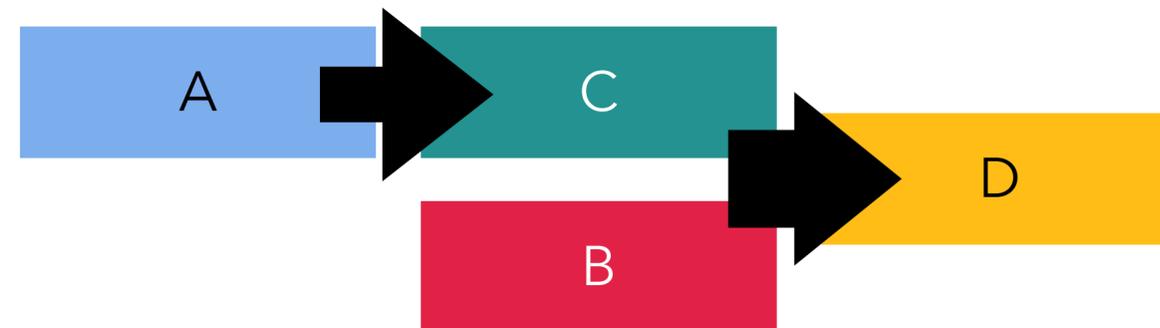
Client-Server



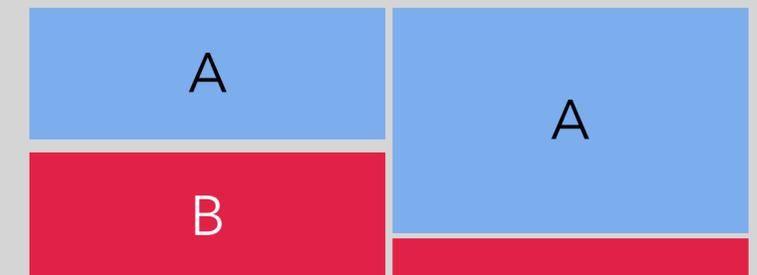
MPMD



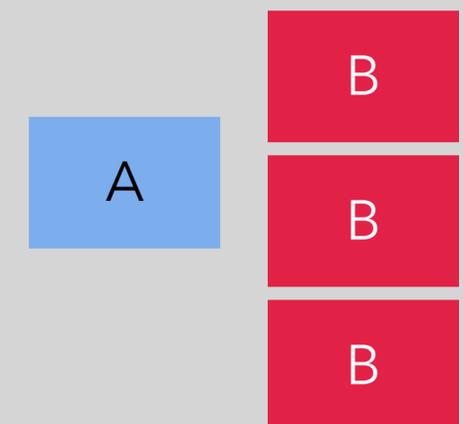
Systolic Architectures / Workflows



Malleability



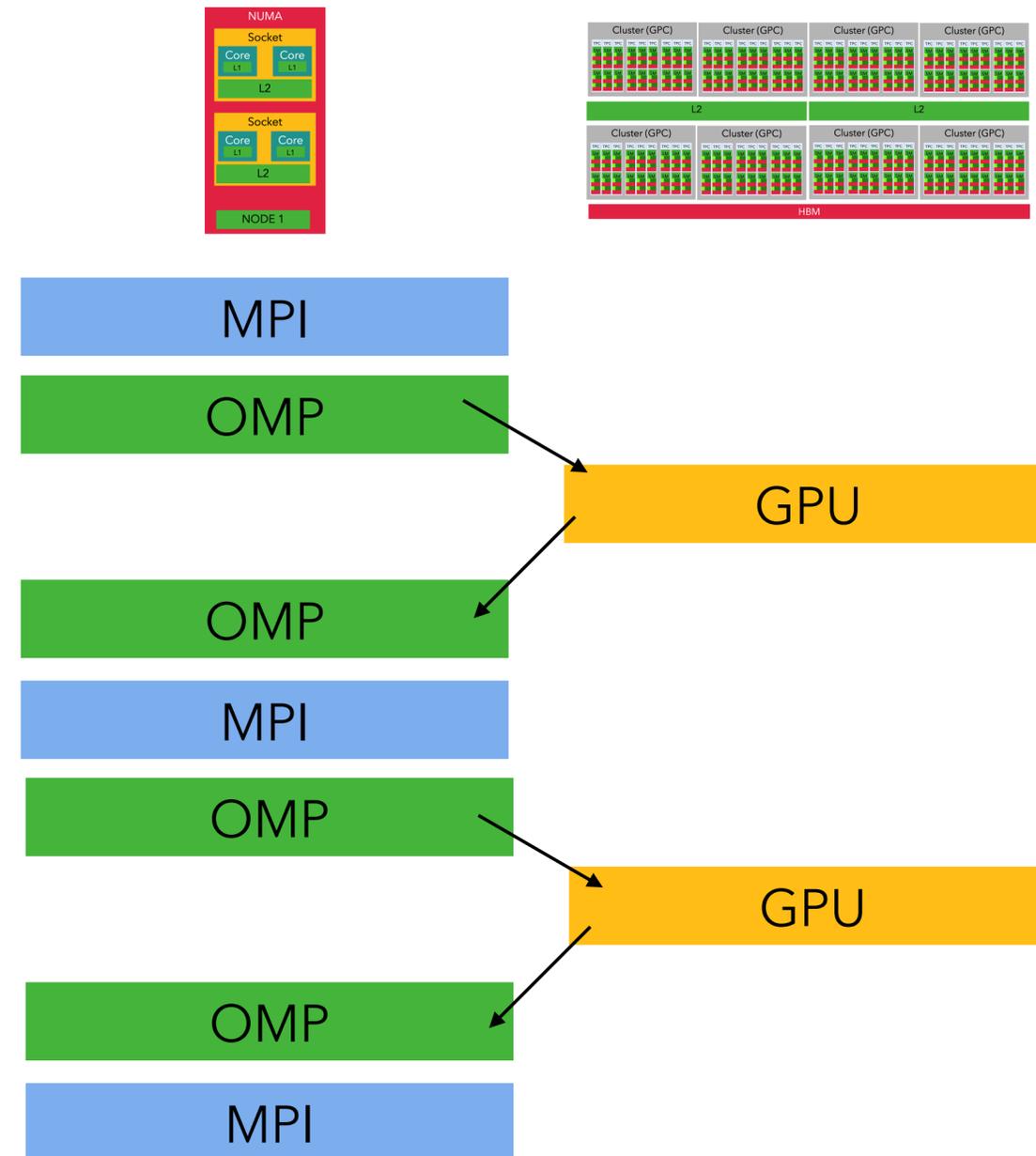
Invasive Computing



On Software Differentiation

CURRENT NESTING

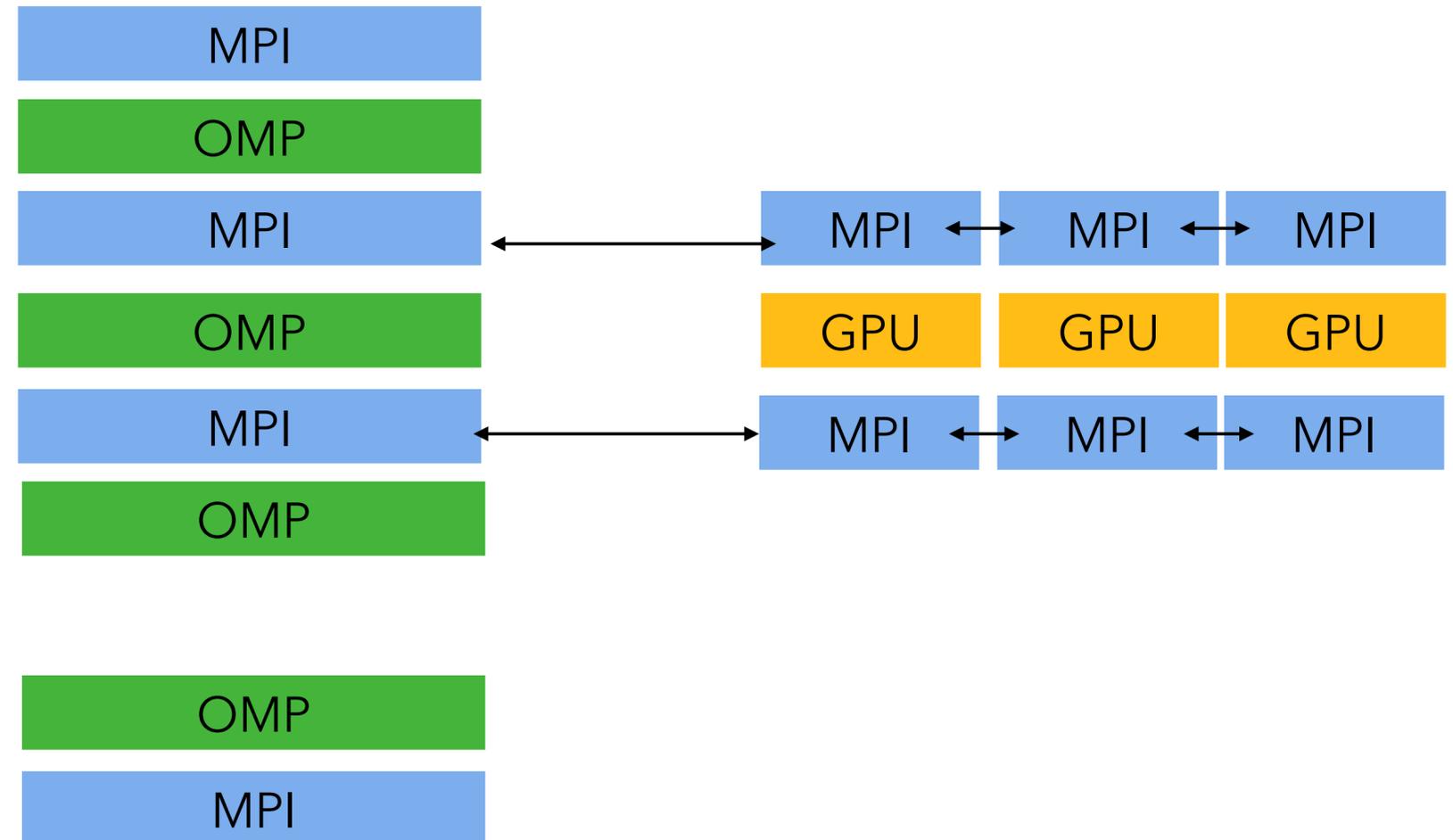
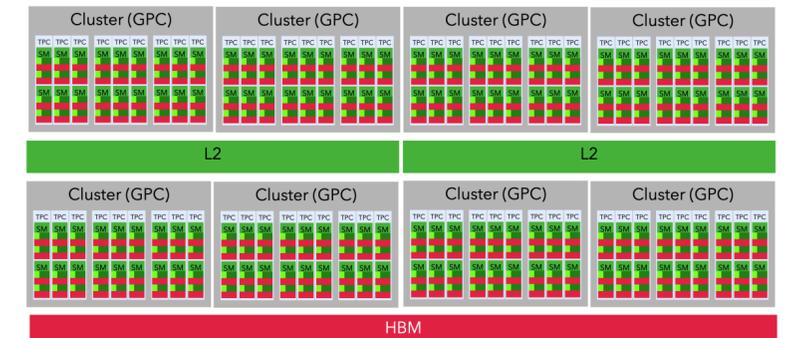
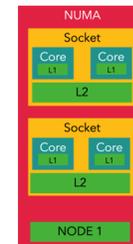
- Alternating behavior
- Difficult to light all the silicium at once
- Some of the **mapping** is outside of MPI's view (CUDA memcpy, CUDA kernels)



On Software Differentiation

DIFFERENTIATION

- Functions / Kernels are spatially bound to a given MPI rank
- Already practical with CCL
- Single indexing scheme (MPI rank)
- MPI is the coordinator (more than doing data-movements)



Current Support of MPI in SHMEM

Shared-Memory Support in MPI

- Evolutions of the standard:
 - Shared-memory windows
 - **Finpoints** -> Partitioned Communications
 - Endpoints
 - Support for RPCs (active messages)
 - Continuations
 - Ownership passing
- Patterns of interest:
 - RPCs (for kernel calls and work dispatch)
 - Shared-memory queues (mostly in language at this point)

Shared-Memory Support in Runtimes

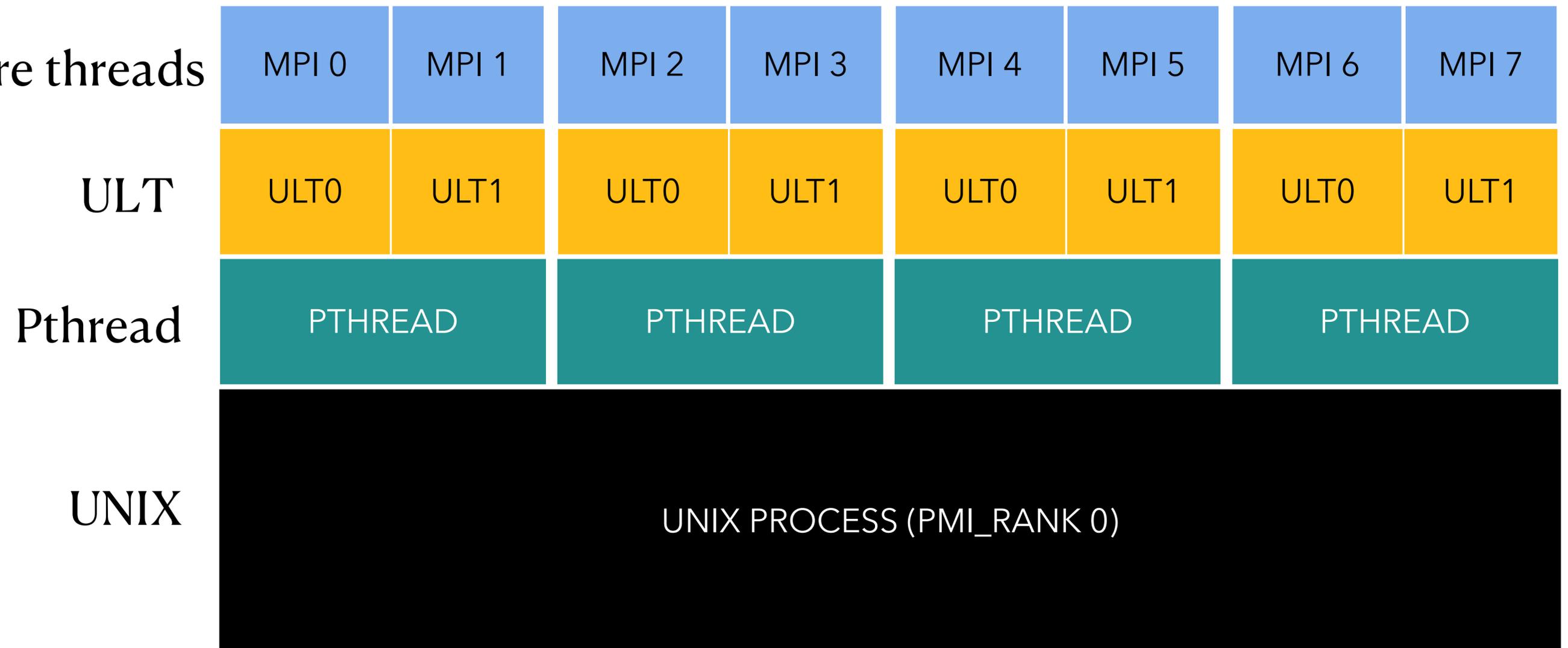
- User-Level:
 - HMPI
 - AMPI
 - Process-In-Process (PiP)
 - Multi-Processor Computing (MPC)
- Kernel level:
 - SMARTMAP
 - PVAS

Shared-Memory Support in Runtimes

- User-Level:
 - HMPI
 - AMPI
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 - **Multi-Processor Computing (MPC)**
- Kernel level:
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 - PVAS

Multi-Processor Computing

MPI Processes are threads



MPC and Existing Codes

```
#include <mpi.h>
#include <stdio.h>

int rank = -1;

int main( int argc, char ** argv )
{
    MPI_Init( &argc, &argv );

    MPI_Comm_rank( MPI_COMM_WORLD, &rank );

    printf("My rank %d", rank );

    MPI_Finalize();

    return 0;
}
```

```
$ mpcrun -p=4 -n=4 ./a.out
```

```
My rank 2
My rank 3
My rank 0
My rank 1
```

It's working with $n == p$ → Process-based MPI

```
$ mpcrun -p=1 -n=4 ./a.out
```

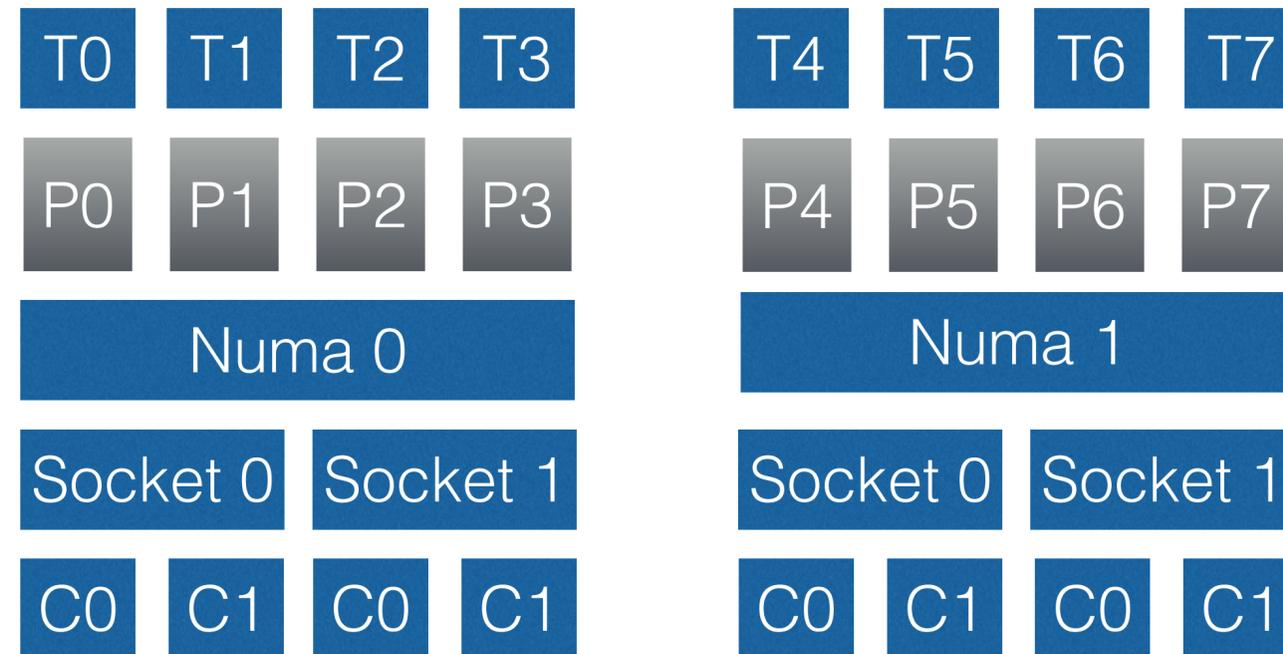
```
My rank 1
My rank 1
My rank 1
My rank 1
```

Overwritten global variable → Thread-based MPI

Need for a privatisation mechanism extending TLS[1]

[1] Jean-Baptiste Besnard, Julien Adam, Sameer Shende, Marc Pérache, Patrick Carribault, Julien Jaeger, and Allen D. Maloney. 2016. Introducing Task-Containers as an Alternative to Runtime-Stacking. In Proceedings of the 23rd European MPI Users' Group Meeting (EuroMPI '16). Association for Computing Machinery, New York, NY, USA, 51–63. <https://doi.org/10.1145/2966884.2966910>

MPC Launch Configuration

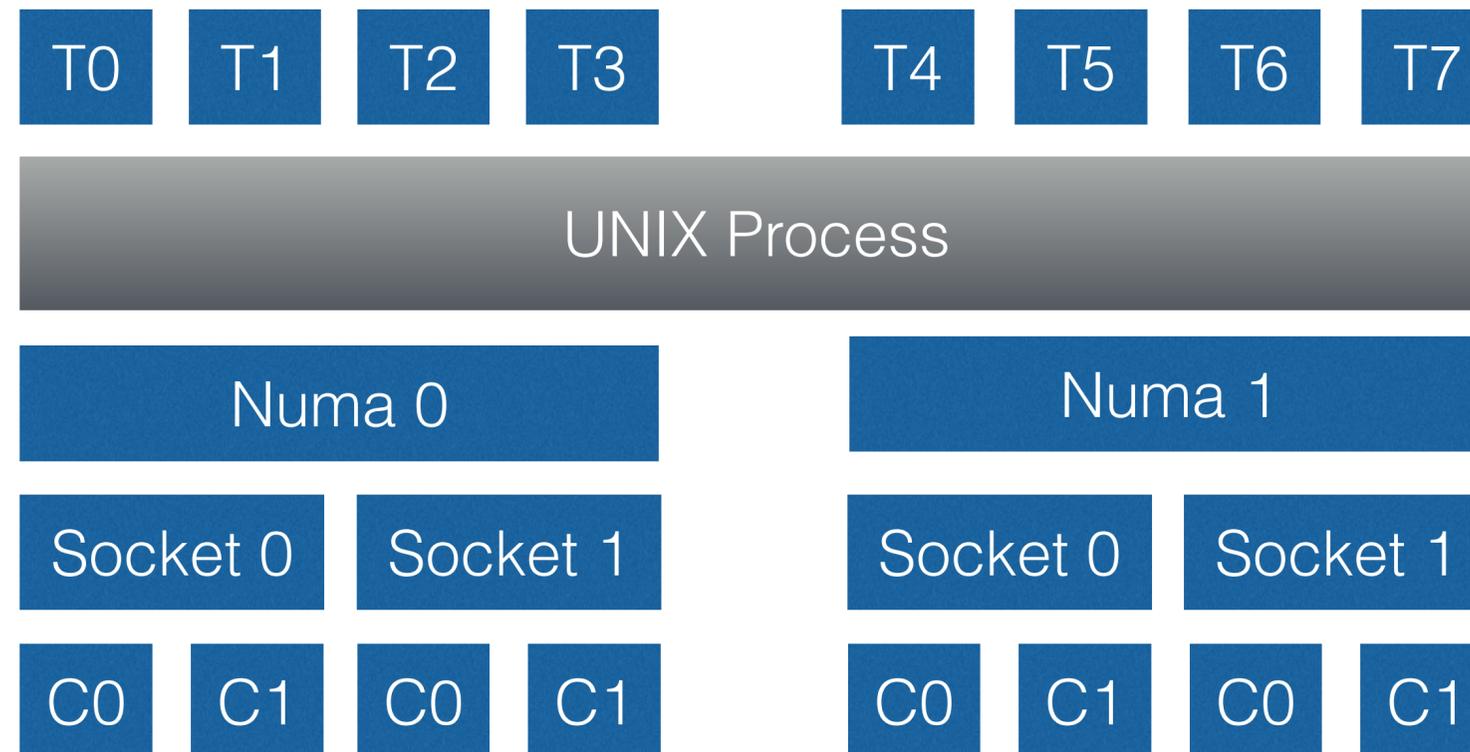


Process-Based MPI
`mpcrun -N=1 -p=8 -n=8 ./a.out`

Launch is described with four parameters (replacing `-np`):

- ▶ **N**: Number of nodes
- ▶ **p**: Number of system processes
- ▶ **n**: Number of MPI tasks
- ▶ **c**: Number of cores per process

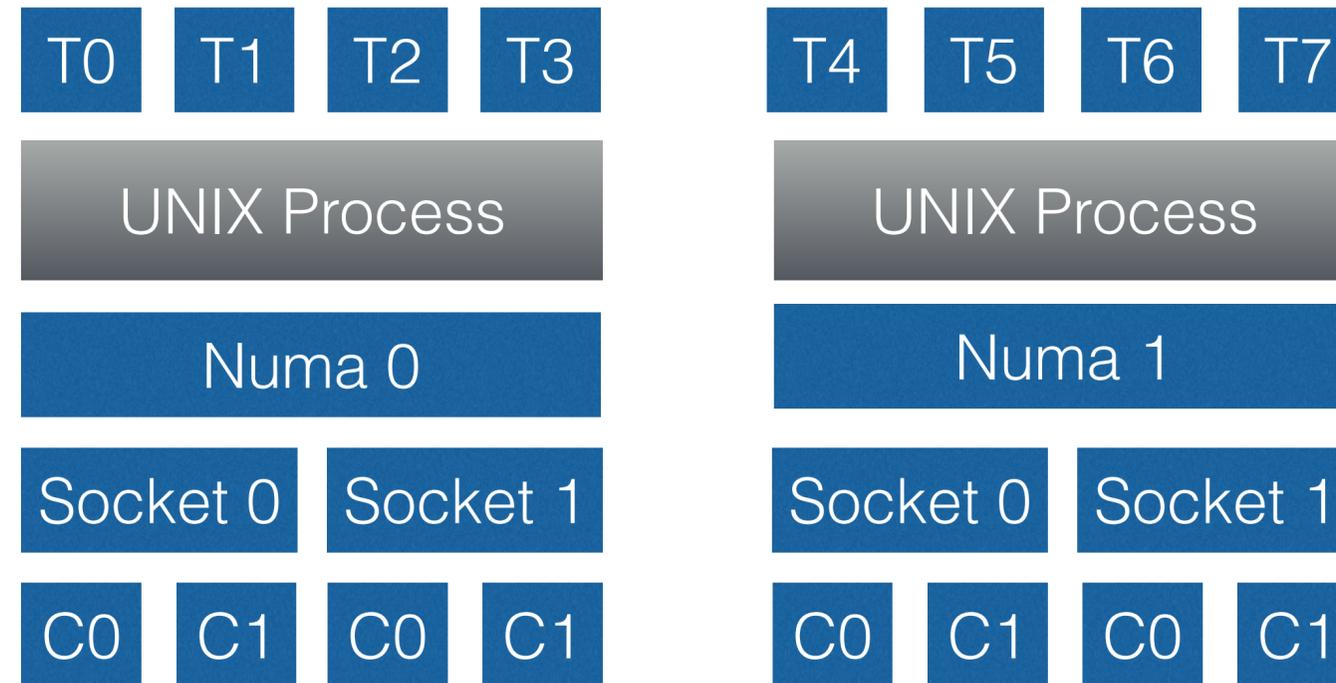
MPC Launch Configuration



Thread-Based MPI

```
mpcrun -N=1 -p=1 -n=8 ./a.out
```

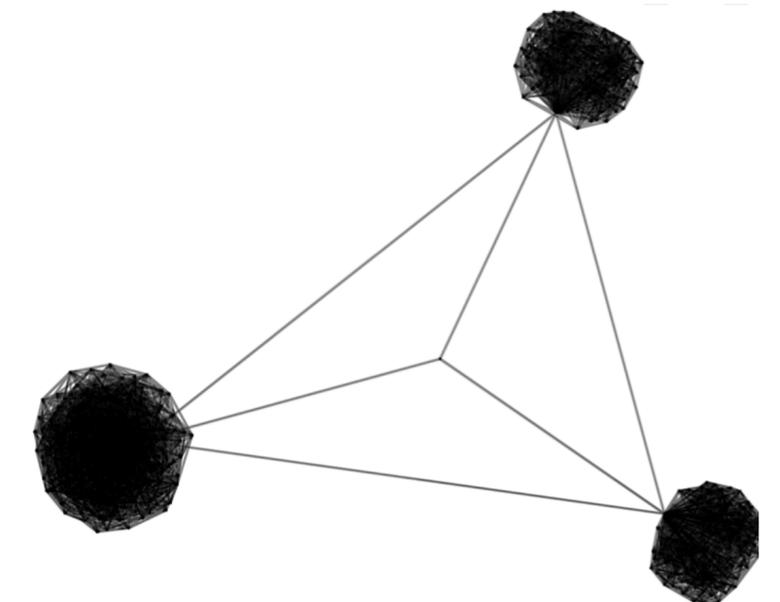
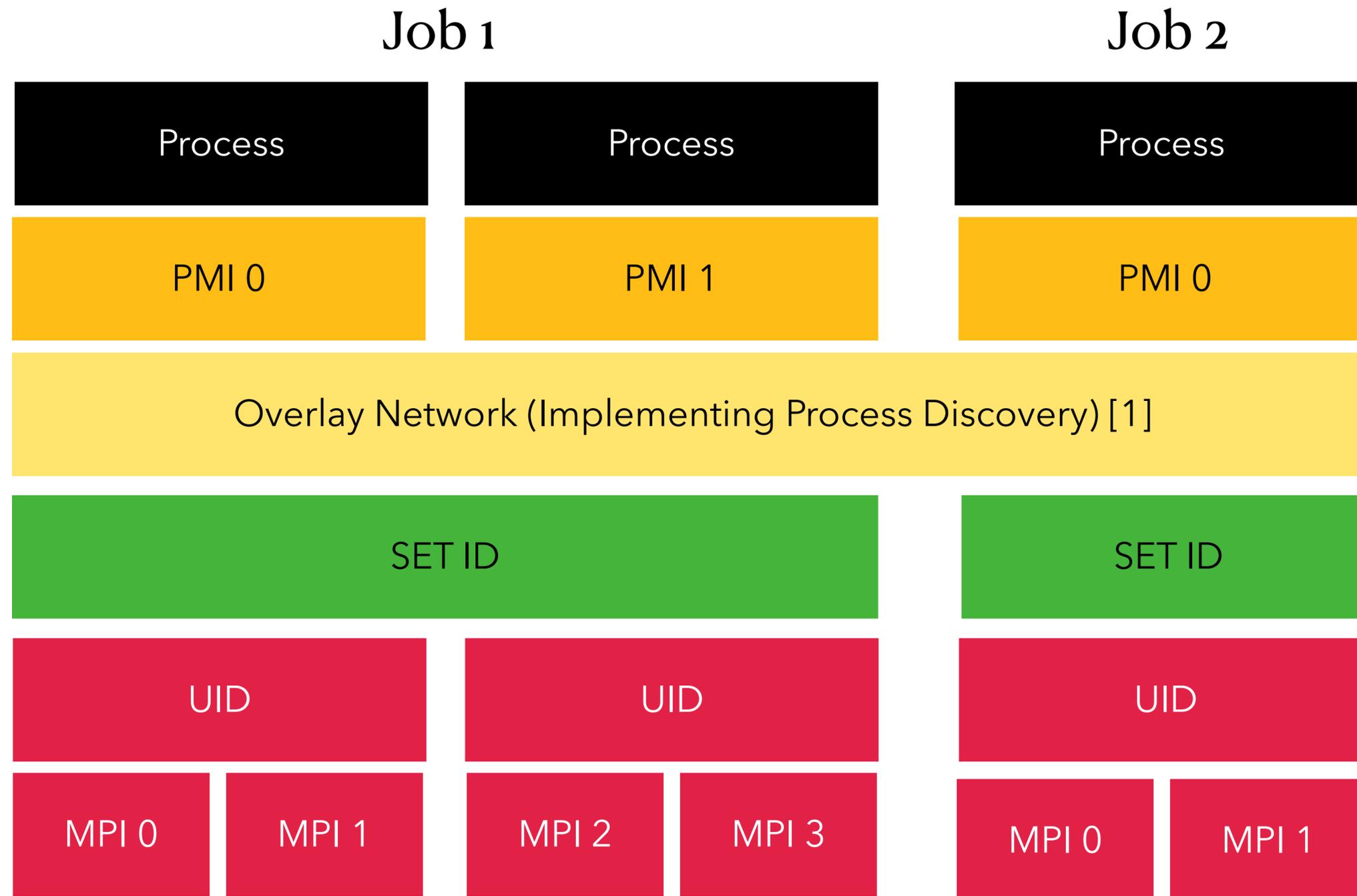
MPC Launch Configuration



Mixed Approach

```
mpcrun -N=1 -p=2 -n=8 -c=4 ./a.out
```

MPC Launch

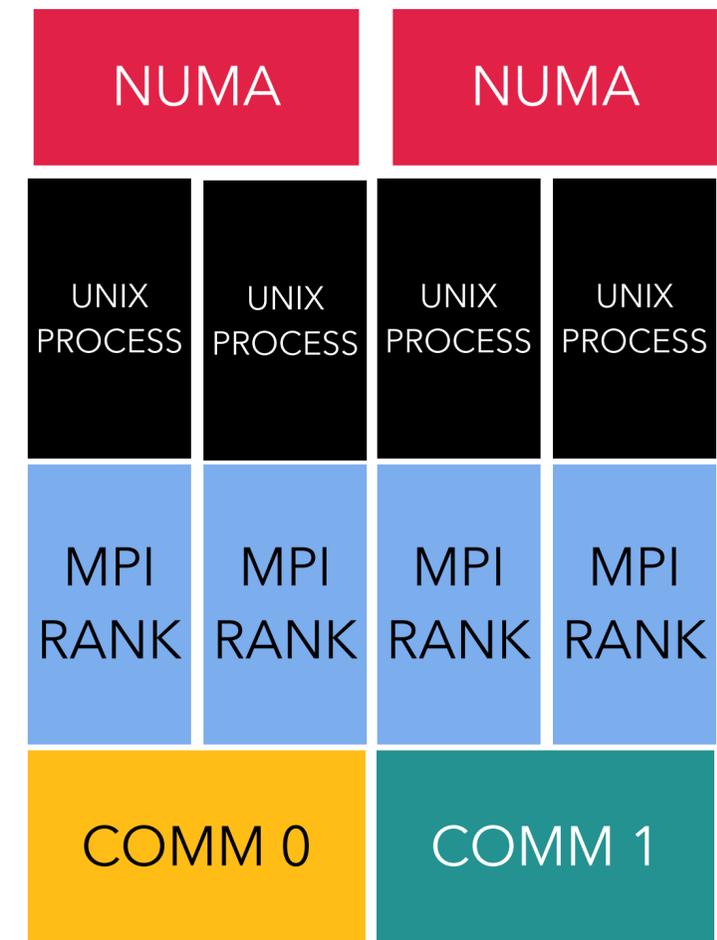
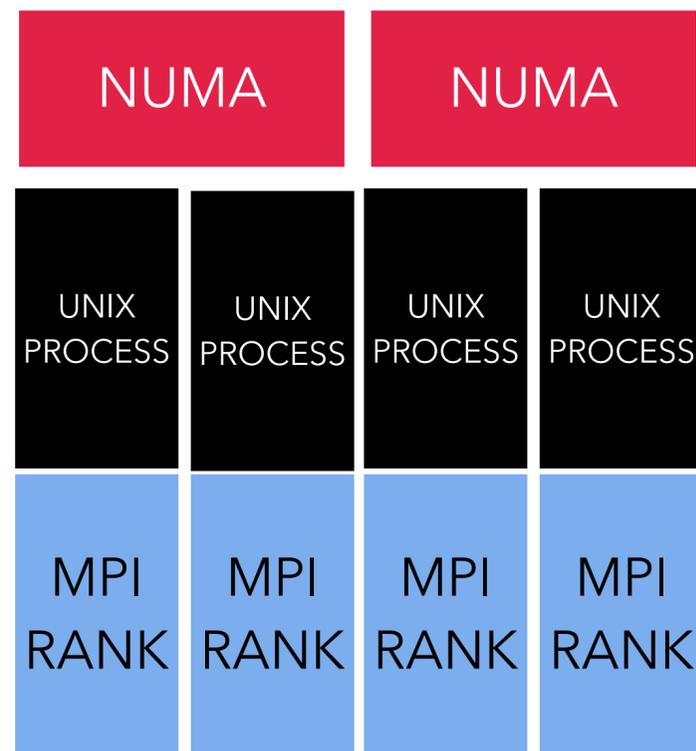


[1] Jean-Baptiste Besnard, Sameer Shende, Allen Malony, Julien Jaeger, and Marc Perache. 2022. Enabling Global MPI Process Addressing in MPI Applications. In Proceedings of the 29th European MPI Users' Group Meeting (EuroMPI/USA '22). Association for Computing Machinery, New York, NY, USA, 27–36. <https://doi.org/10.1145/3555819.3555829>

Sessions at the Rescue

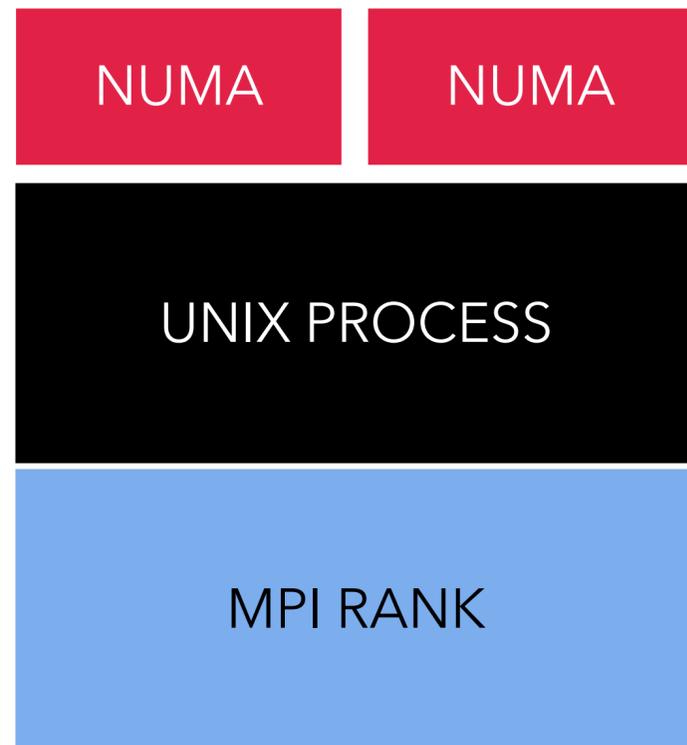
Topological Split on MPI Processes

MPI's role is to **MAP** and **ADDRESS** the distributed MPI processes



Topological Split Inside Processes

WORLD MODEL

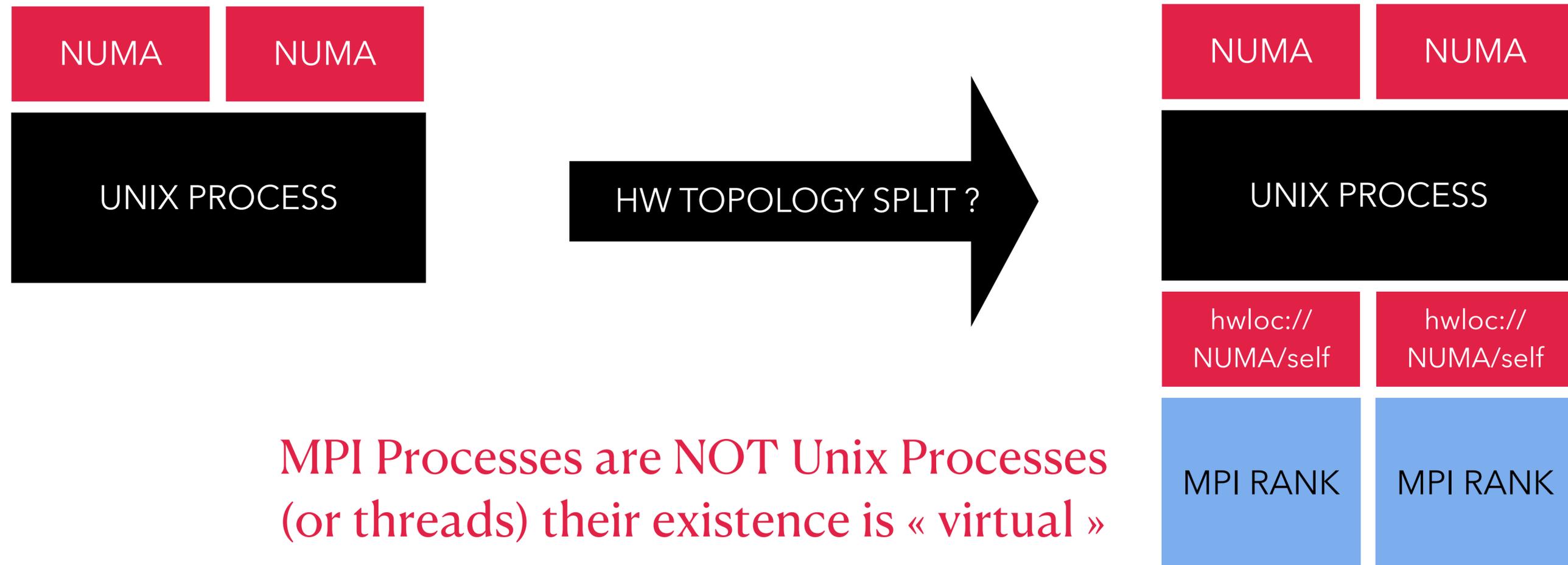


?

Rank is bound to proc

Topological Split Inside Processes

SESSIONS MODEL



MPI Processes are NOT Unix Processes
(or threads) their existence is « virtual »

Context inherited by handle nesting

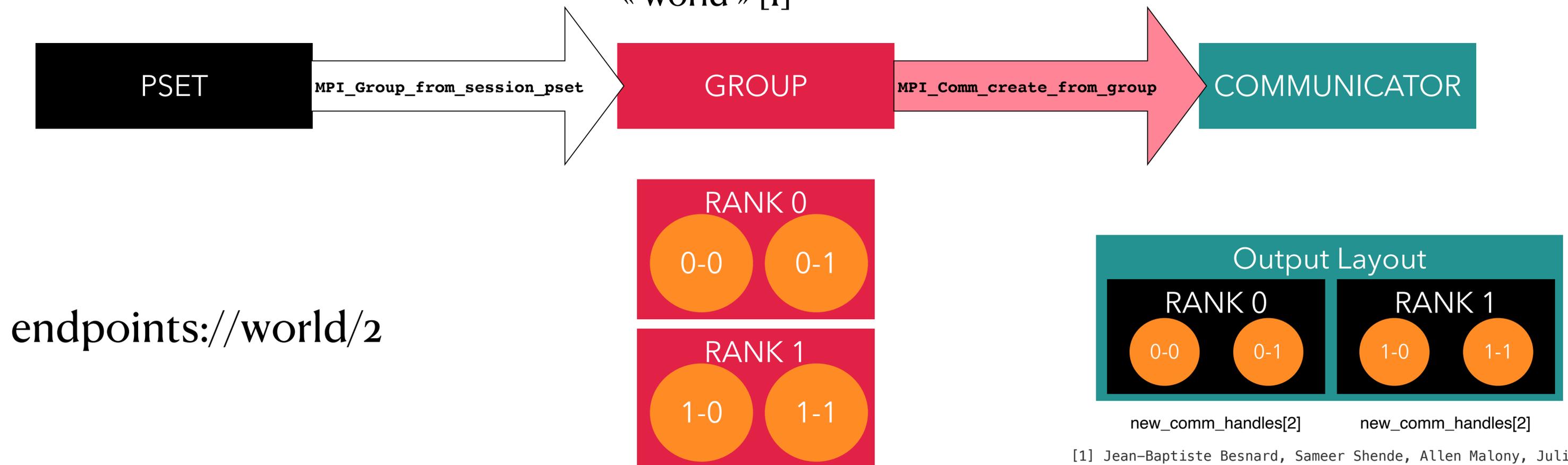
Topological Split Inside Processes

SESSIONS MODEL

Groups with processes that are initially not from « world » [1]

Instead, communicators created in a non-collective manner (not bound to executors) == **Endpoints**

```
int MPI_Comm_endpoints_from_group(MPI_Group group,  
MPI_Info info, MPI_Comm *new_comm_handles,  
int max_len, char *tag)
```



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